Compiler architecture: the discipline and abstractions to exploit application domains to target accelerators to feed microarchitecture

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Joint work with Luke Panayi, Martin Berger, Rohan Gandhi, Jim Whittaker, Vassilios Chouliaras, Andrei Sburlan, Avaneesh Deleep, George Bisbas, Edward Stow, Jacky Wong, David Ham, Fabio Luporini, Lawrence Mitchell, Graham Markall, Mike Giles, Gerard Gorman, Florian Rathgeber, Luigi Nardi, Carlos Cueto, Lluis Guasch, Fabio Luporini, Oscar Bates, George Strong, Oscar Calderon Agudo, Javier Cudeiro, Gerard Gorman. and Meng-Xing Tang, Hidenobu Matsuki, Riku Murai, Andrew J. Davison, David Pearce, Piotr Dudek, Tobias Grosser and many more

Who am I and what do I do?

- I've worked on a lot of things....
 - GPUs, FPGAs, cache coherency, vectorisation, verification, bounds checking, CPU architecture...
- I have worked on general-purpose compilers
 - Notably pointer analysis
 - adopted into GCC
 - (actually the work of my PhD graduate David Pearce)
- But the benefits were incremental

- Meanwhile I engaged with applications specialists
 - Who know they have major performance optimisation opportunities
 - So I got interested in automating domain-specific optimisations
- Which grew into deep engagements with projects like
 - Firedrake, Devito and PyFR, that automate the pathway from PDE to high-performance code for laptops and supercomputers
 - Robot vision, SLAM, robot localisation

Vectorisation, parametric polyhedral tiling

Tiling for unstructured-mesh stencils

Lazy, data-driven computecommunicate

Runtime code generation

Multicore graph worklists

Tensor contractions

Generalised loop-invariant code motion

Sparsity in Fourier transforms

Functional Variational Inference

Search-based optimisation

Processor/accelerator microarchitecture, co-design

MLIR

Technologies

Finite-volume CFD

Finite-element

Finite-difference

Real-time 3D scene understanding

Adaptive-mesh CFD

Contour trees, Reeb graphs

Unsteady CFD - higherorder flux-reconstruction

Ab-initio computational chemistry (ONETEP)

Gaussian belief propagation

Uncertainty in DNNs

Near-camera processing

Quantum computing

Contexts

PyOP2/OP2

Unstructured-mesh stencils

Firedrake

Automating

performance

optimisations

Exploiting

higher-level

language to

performance

get better

than low

level code

domain-

specific

Finite-element

Devito: finite difference

SLAMBench:

3D vision, dense SLAM

SuperEight: octree SLAM

PRAgMaTIc: Unstructured mesh adaptation

GiMMiK: small matrix multiply

TINTL: Fourier interpolation

Hypermapper: design optimisation

RobotWeb: distributed localisation

CAIN: convolutions on analogue SIMD focal-plane processor

Projects

Aeroengine turbomachinery

Weather and climate

Glaciers

Domestic robotics, augmented reality

Tidal turbine placement

Formula-1, UAVs, buildings

Solar energy, drug design

Energy storage

Medical imaging

Remote sensing in robot agriculture

Applications

What are we doing next?

- Gaussian Belief Propagation as a foundation for managing locality, distribution, asynchrony and approximation in spatial Al
- Gaussian Splatting SLAM: capturing manipulable photorealistic scenes in real time
- Quantum circuits as a DSL

- XDSL and MLIR: common compiler architecture ecosystem for DSLs
- Instruction labelling: policy/prediction hints for scalable microarchitecture
- Tensor contractions as a compiler IR
- On-sensor vision computing/accelerator architectures and programming models
- Compiler technology for simulating quantum computers

This talk

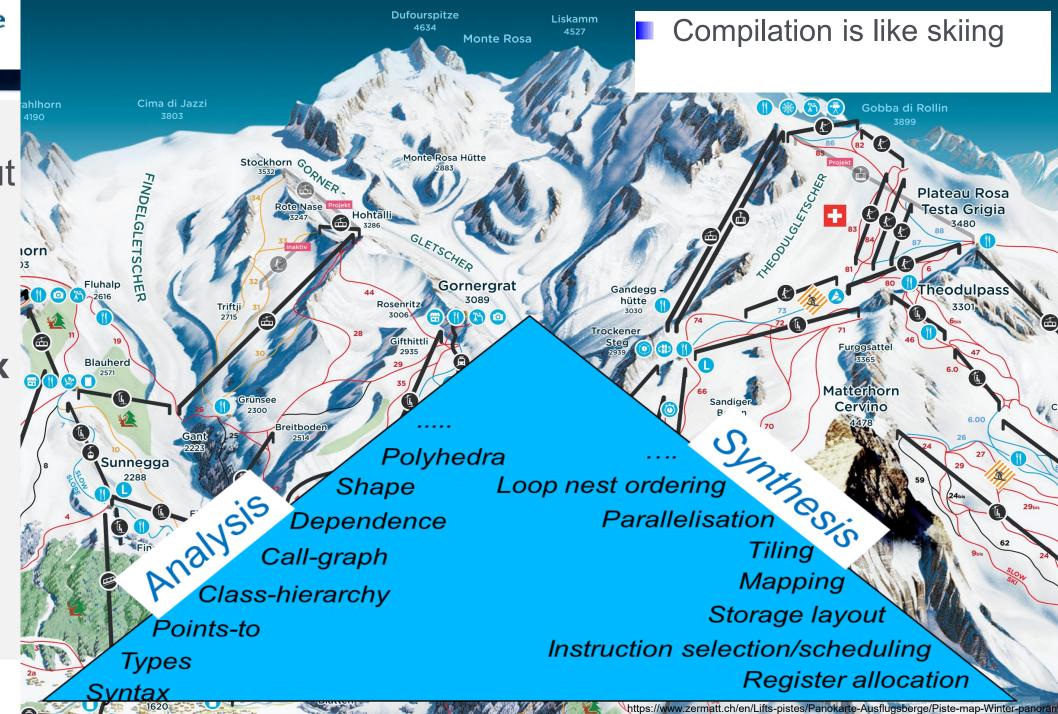
- Compiler architecture
- The importance of capturing application code at the highest-possible level of representation
- Getting the abstraction right
 - To capture what the code is trying to do
 - To capture what we need to optimise for the hardware
 - For CPUs
 - **For GPUs**
 - **■** For custom accelerators

Turing tax

- Alan Turing realised we could use digital technology to implement any computable function
- He then proposed the idea of a "universal" computing device a *single* device which, with the right program, can implement any computable function *without further configuration*
- "Turing Tax", or "Turing Tariffs": the overhead (performance, cost, or energy) of universality in this sense
- The performance (time/area/energy) difference between a special-purpose device and a general-purpose one
- One of the fundamental questions of computer architecture is to how to reduce the Turing Tax

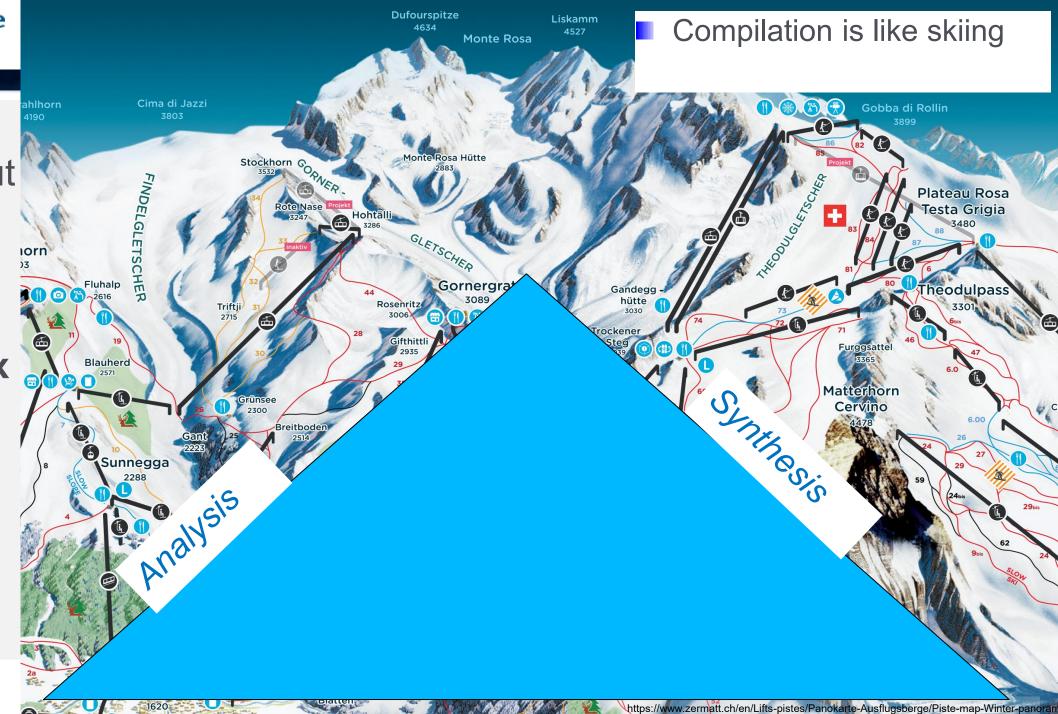
What about compilers ?

Is there a Turing Tax for compilers too?



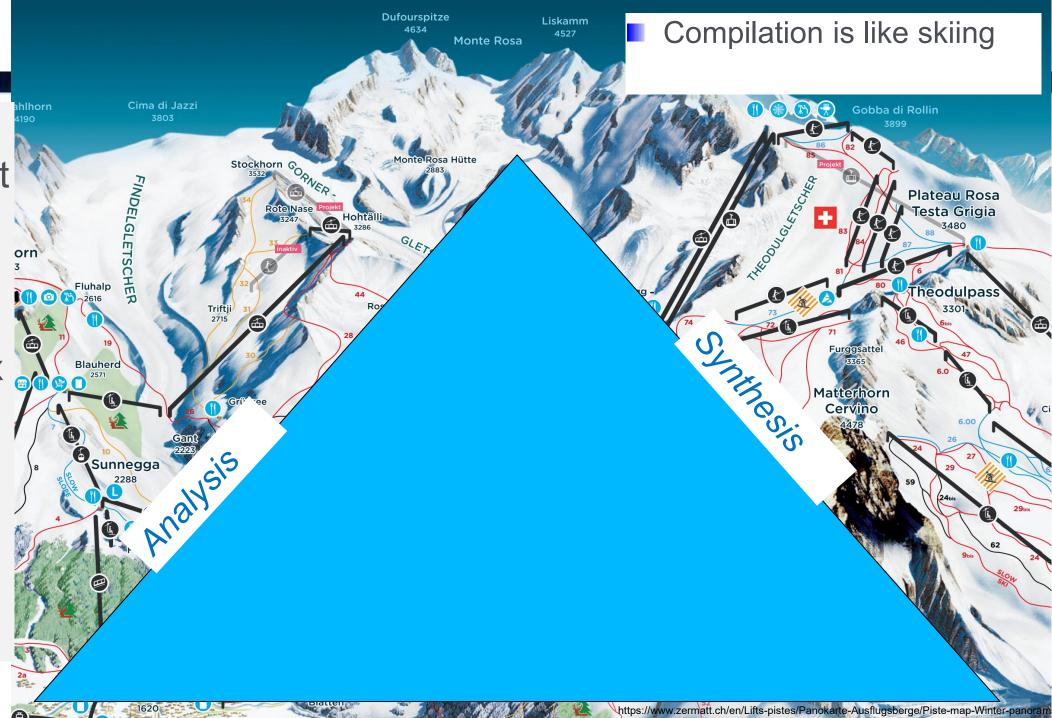
What about compilers?

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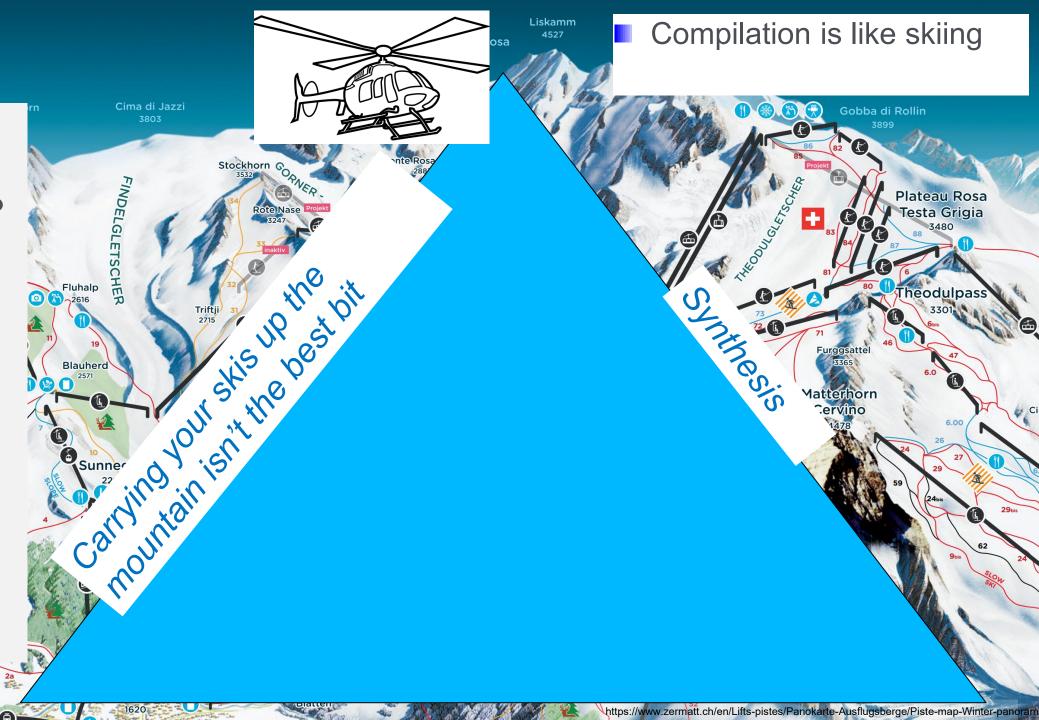


What about compilers?

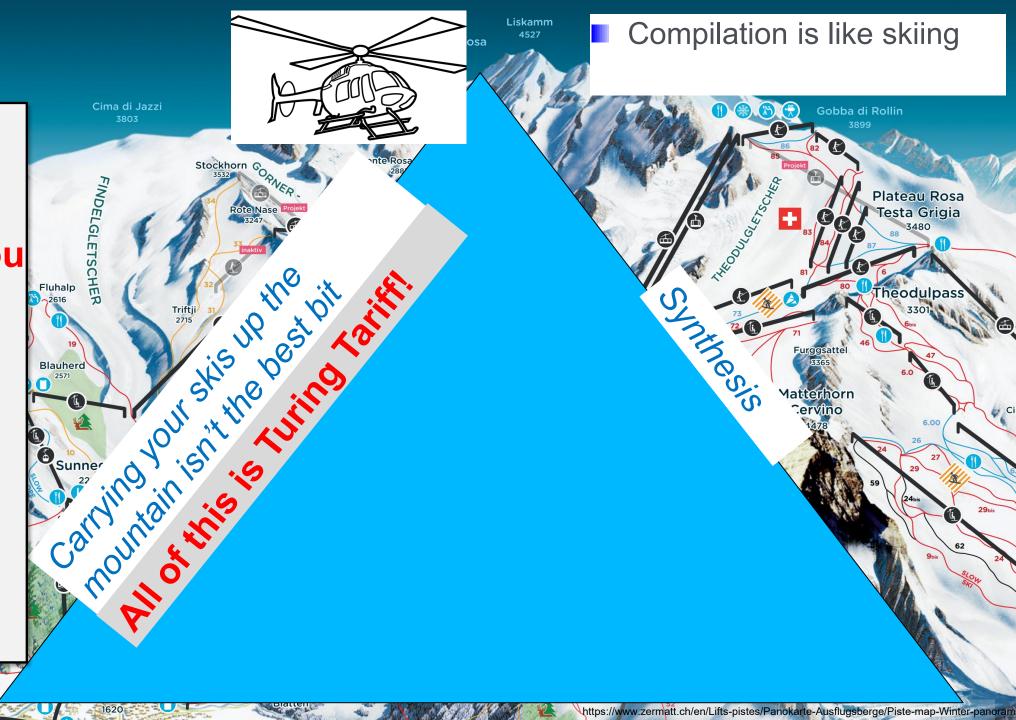
Is there a
Turing Tax
for
compilers
too?



- What about compilers?
- Is there a
 Turing Tax
 for
 compilers
 too?



- What about compilers?
- The price you pay for coding in a general-purpose language
- When you could have used a DSL





-trearak

Documentation

Install

Citing

Team

Publications

Events

Funding

Contact

GitHub

















Features:

Expressive specification of any PDE using the Unified Form Language from the FEniCS Project.

Firedrake is an automated system for the solution of partial differential equations

using the finite element method (FEM). Firedrake uses sophisticated code gener-

ation to provide mathematicians, scientists, and engineers with a very high pro-

ductivity way to create sophisticated high performance simulations.

- Sophisticated, programmable solvers through seamless coupling with PETSc.
- Triangular, quadrilateral, and tetrahedral unstructured meshes.
- Layered meshes of triangular wedges or hexahedra.
- Vast range of finite element spaces.
- Sophisticated automatic optimisation, including sum factorisation for high order elements, and vectorisation.
- Geometric multigrid.
- Customisable operator preconditioners.
- Support for static condensation, hybridisation, and HDG methods.

Latest commits to the Firedrake main branch on **Github**

Merge pull request #4536 from firedrakeproject/JHopeCollins/mergerelease-into-main

Connor Ward authored at 05/09/2025. 15:47:46

Merge branch 'main' into JHopeCollins/merge-release-into-main Josh Hope-Collins authored at 05/09/2025. 14:34:12

Fix docs build (#4537) Connor Ward authored at 05/09/2025. 14:31:29

MG: rediscretize with a different coarse_mat_type (#4538)
Pablo Brubeck authored at 05/09/2025, 10:35:14

Remove `Interpolator.interpolate` (#4531) Leo Collins authored at 04/09/2025, 16:34:07

ormer team members





























Firedrake/FEniCS ecosystem

- Many projects build on Firedrake, for example:
 - Firedrake-adjoint: extension of the pyadjoint algorithmic differentiation framework to yield fully automated derivation of adjoint PDE solvers
 - Irksome: Automates Runge–Kutta time-stepping
 - AsQ: parallel-in-time
 - Slate: linear algebra on finite element tensors, for hybridisation, static condensation
 - PCPATCH: topological construction of multigrid relaxation methods

- Defcon: deflated continuation method for computing bifurcation diagrams
- Goal-oriented mesh adaptation
- Integration with PyTorch and JAX

Including tools for specific application domains, for example:

- Thetis: unstructured grid coastal modelling framework
- IcePack: glacier flow
- Gusto: dynamical cores for weather prediction and climate models



Firedrake example: Burgers equation

$$\int_{\Omega} \frac{u^{n+1} - u^n}{dt} \cdot v + ((u^{n+1} \cdot \nabla)u^{n+1}) \cdot v + \nu \nabla u^{n+1} \cdot \nabla v \, \mathrm{d}x = 0.$$

- From the weak form of the PDE, we derive an equation to solve, that determines the state at each timestep in terms of the previous timestep
- Transcribe into Python: u is u^{n+1} , u_ is u^n :

 helicopter ride

Set up the equation and solve for the next timestep u:

$$solve(F == 0, u)$$

At this point, Firedrake generates code to assemble a linear system, runs it and calls a linear solver (we use PetSC)

```
from firedrake import *
n = 50
mesh = UnitSquareMesh(n, n)
# We choose degree 2 continuous Lagrange polynomials. We also need a
# piecewise linear space for output purposes::
V = VectorFunctionSpace(mesh, "CG", 2)
V out = VectorFunctionSpace(mesh, "CG", 1)
# We also need solution functions for the current and the next timestep::
u = Function(V, name="Velocity")
u = Function(V, name="VelocityNext")
v = TestFunction(V)
# We supply an initial condition::
x = SpatialCoordinate(mesh)
ic = project(as vector([sin(pi*x[0]), 0]), V)
# Start with current value of u set to the initial condition, and use the
# initial condition as our starting guess for the next value of u::
u .assign(ic)
u.assign(ic)
# :math:`\nu` is set to a (fairly arbitrary) small constant value::
nu = 0.0001
timestep = 1.0/n
# Define the residual of the equation::
F = (inner((u - u)/timestep, v))
     + inner(dot(u,nabla_grad(u)), v) + nu*inner(grad(u), grad(v)))*dx
outfile = File("burgers.pvd")
outfile.write(project(u, V_out, name="Velocity"))
# Finally, we loop over the timesteps solving the equation each time::
t = 0.0
end = 0.5
while (t <= end):</pre>
    solve(F == 0, u)
    u .assign(u)
    t += timestep
    outfile.write(project(u, V out, name="Velocity"))
```

Burgers equation

- Complete runnable Python script
 - sets up the equation,
 - specifies discretisation
 - Iterates multiple timesteps
 - Writes velocity field out for animated visualisation
- Firedrake implements the Unified Form Language (UFL)
- Embedded in Python



(UFL is also the DSL of the FEniCS project)

```
mesh = UnitSquareMesh(n, n)
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u = Function(V, name="VelocityNext")
                                                            = Function(V, name="VelocityNext")
v = TestFunction(V)
# We supply an initial condition::
                                                          # set up initial conditions for u and u
x = SpatialCoordinate(mesh)
ic = project(as vector([sin(pi*x[0]), 0]), V)
# Start with current value of u set to the initial condition, and use the
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timestep = 1.0/n
                                                               + inner(dot(u,nabla_grad(u)), v) + nu*inner(grad(u), grad(v)))*dx
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                                                          t = 0.0
   + inner(dot(u,nabla grad(u)), v) + nu*inner(grad(u), grad(v)))*dx
                                                                                               Time-stepping loop:
                                                         end = 0.5
outfile = File("burgers.pvd")

    Solve for the state at the next

                                                         while (t <= end):</pre>
outfile.write(project(u, V out, name="Velocity"))
                                                                                                  timestep
# Finally, we loop over the timesteps solving the equation each time::
                                                                solve(F == 0, u)
t = 0.0
                                                                                                  Save snapshot to file for
                                                                u_.assign(u)
end = 0.5
while (t <= end):
                                                                                                  animation
                                                                t += timestep
   solve(F == 0, u)
   u_.assign(u)
   t += timestep
                                                                outfile.write(project(u, V_out, name="Velocity"))
   outfile.write(project(u, V out, name="Velocity"))
```

from firedrake import *

Gather data from neighbouring cells using adjacency graph

Compute re-used common sub-terms

Compute integrals using polynomial approximation of fluid state in this cell

Sum local contributions into global system matrix

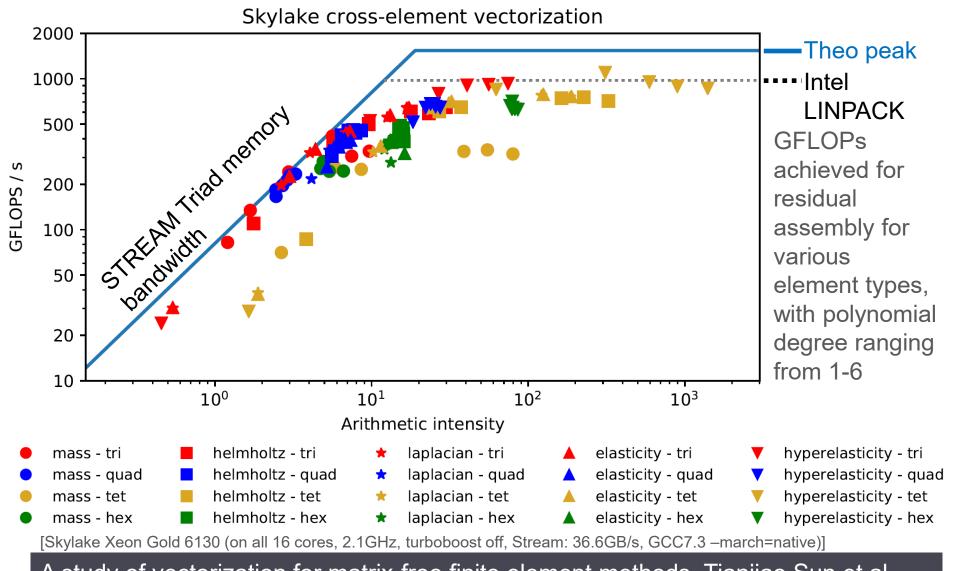
```
#include <math.h>
void wrap form00 cell integral otherwise(int const start, int const end, Mat const mat0, double const * restrict dat1, double const * restrict dat0, int const * restrict map0, int const * restrict map0, int const * restrict map1)
 double const form t17[7] = { ... };
double const form_t18[7 * 6] = { ... };
double const form_t19[7 * 6] = { ... };
  double form_t2;
  double const form_t20[7 * 6] = { ... };
                                                                                                                                                                                                                  Generated code
   double form_t21...t37;
   double form_t38[6];
   double form_t39[6];
   double form_t4;
  double form_t40...t45;
   double form t5...t9:
                                                                                                                                                                                                                                  to assemble the
 double t0[6 * 2];
double t1[3 * 2];
double t2[6 * 2 * 6 * 2];
   for (int n = start; n <= -1 + end; ++n)
                                                                                                                                                                                                                                   resulting linear
     for (int i4 = 0; i4 <= 5; ++i4)
        for (int i5 = 0; i5 \le 1; ++i5)
           for (int i6 = 0; i6 <= 5; ++i6)
              for (int i7 = 0; i7 <= 1; ++i7)

t2[24 * i4 + 12 * i5 + 2 * i6 + i7] = 0.0;
                                                                                                                                                                                                                                  system matrix
     for (int i3 = 0; i3 <= 1; +i3)

for (int i3 = 0; i3 <= 1; +i3)

t1[2 * i2 + i3] = dat1[2 * map1[3 * n + i2] + i3];
    for (int i0 = 0; i0 < = 5; ++10)
for (int i1 = 0; i0 < = 5; ++10)
for (int i1 = 0; i1 <= 1; ++i1)
t0[2 * i0 + i1] = dat0[2 * map0[6 * n + i0] + i1];
form_t0 = -1.0 * t1[1];</pre>
                                                                                                                                                                                                                  Executed at each
     form_t1 = form_t0 + t1[3];
form t2 = -1.0 * t1[0];
     form_t3 = form_t2 + t1[2];
form t4 = form t0 + t1[5]:
                                                                                                                                                                                                                                  triangle in the
     form_t5 = form_t2 + t1[4];
form t6 = form_t3 * form_t4 + -1.0 * form_t5 * form_t1;
     form_t7 = 1.0 7 form_t6;
form_t8 = form_t7 * -1.0 * form_t1;
form_t9 = form_t4 * form_t7;
                                                                                                                                                                                                                                   mesh
      form t10 = form t3 * form t7;
      form t11 = form t7 * -1.0 * form t5:
      form t12 = 0.0001 * (form t8 * form t9 + form t10 * form t11);
      form t13 = 0.0001 * (form t8 * form t8 + form t10 * form t10);
      form t14 = 0.0001 * (form t9 * form t9 + form t11 * form t11);
                                                                                                                                                                                                                             Accesses
      form_t15 = 0.0001 * (form_t9 * form_t8 + form_t11 * form_t10);
      form t16 = fabs(form t6);
      for (int form_ip = 0; form_ip <= 6; ++form ip)
         form t26 = 0.0; form t25 = 0.0; form t24 = 0.0; form t23 = 0.0; form t22 = 0.0; form t21 = 0.0;
                                                                                                                                                                                                                                   degrees of
          for (int form_i = 0; form_i <= 5; ++form_i)
            form t21 = form \ t21 + form \ t20[6 * form ip + form i] * t0[1 + 2 * form i];
            form t22 = form t22 + form t19[6 * form ip + form i] * t0[1 + 2 * form i];
            form t23 = form t23 + form t20[6 * form ip + form i] * t0[2 * form i];
                                                                                                                                                                                                                                   freedom shared
            form t24 = form t24 + form t19[6 * form ip + form i] * t0[2 * form i];
            form t25 = form t25 + form t18[6 * form ip + form i] * t0[1 + 2 * form i];
            form_t26 = form_t26 + form_t18[6 * form_ip + form_i] * t0[2 * form_i];
          form_t27 = form_t17[form_ip] * form_t16;
          form t28 = form t27 * form t15;
                                                                                                                                                                                                                                  with neighbour
          form t29 = form t27 * form t14;
          form_t30 = form_t27 * (form_t26 * form_t9 + form_t25 * form_t11);
          form_t31 = form_t27 * form_t13;
          form_t32 = form_t27 * form_t12;
          form_t33 = form_t27 * (form_t26 * form_t8 + form_t25 * form_t10);
                                                                                                                                                                                                                                  triangles through
          form t34 = form t27 * (form t11 * form t24 + form t10 * form t23);
          form_t35 = form_t27 * (form_t9 * form_t22 + form_t8 * form_t21);
         form_t36 = form_t27 * (50.0 + form_t9 * form_t24 + form_t8 * form_t23);
         form t37 = form t27 * (50.0 + form t11 * form t22 + form t10 * form t21); for (int form k\theta = 0; form k\theta <= 5; ++form k\theta)
                                                                                                                                                                                                                                   indirection map
            form_t38[form_k0] = form_t18[6 * form_ip + form_k0] * form_t37;
            form_t39[form_k0] = form_t18[6 * form_ip + form_k0] * form_t36;
          for (int form_j0 = 0; form_j0 <= 5; ++form_j0;
           form_t40 = form_t18[6 * form_ip + form_j0] * form_t35; form_t41 = form_t18[6 * form_ip + form_j0] * form_t34; form_t42 = form_t18[6 * form_ip + form_j0] * form_t31 + form_t18[6 * form_ip + form_j0] * form_t33 + form_t19[6 * form_ip + form_j0] * form_t32; form_t43 = form_t26[6 * form_ip + form_j0] * form_t28 + form_t18[6 * form_ip + form_j0] * form_t28 + form_t29; form_t28 + form_t28
            for (int form k0_0 = 0; form k0_0 <= 5; ++form k0_0)
              (form_t44 = form_t43 * form_t19[6 * form_ip + form_k0_0];
form_t45 = form_t42 * form_t20[6 * form_ip + form_k0_0];
form_t45 = form_t42 * form_t20[6 * form_ip + form_k0_0];
t2[24 * form_j0 + 2 * form_k0_0] = t2[24 * form_j0 + 2 * form_k0_0] + form_t45 + form_t18[6 * form_ip + form_j0] * form_t30[form_k0_0] + form_t44;
t2[13 + 24 * form_j0 + 2 * form_k0_0] = t2[13 + 24 * form_j0 + 2 * form_k0_0] + form_t45 + form_t18[6 * form_ip + form_j0] * form_t38[form_k0_0] + form_t44;
t2[1 + 24 * form_j0 + 2 * form_k0_0] = t2[12 + 24 * form_j0 + 2 * form_k0_0] + form_t18[6 * form_ip + form_k0_0] * form_t40;
t2[12 + 24 * form_j0 + 2 * form_k0_0] = t2[12 + 24 * form_j0 + 2 * form_k0_0] + form_t18[6 * form_ip + form_k0_0] * form_t40;
      MatSetValuesBlockedLocal(mat0, 6, &(map0[6 * n]), 6, &(map0[6 * n]), &(t2[0]), ADD VALUES);
```

Does it generate good code?

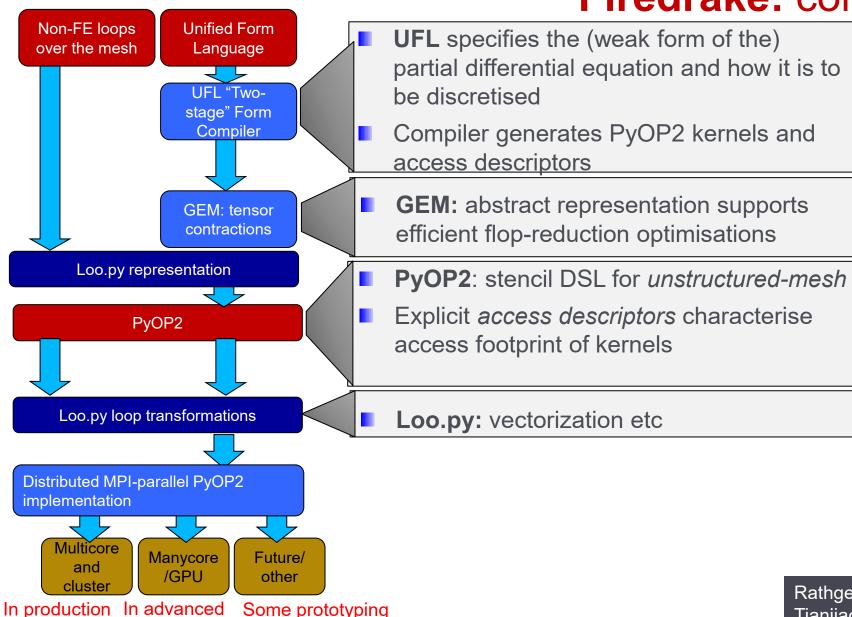


The big story here:

- Five different PDEs
- Six different choices for polynomial degree
- Automatically compiled
- So we see how performance varies across the benchmark suite

A study of vectorization for matrix-free finite element methods, Tianjiao Sun et al IJHPCA 2020 https://arxiv.org/abs/1903.08243

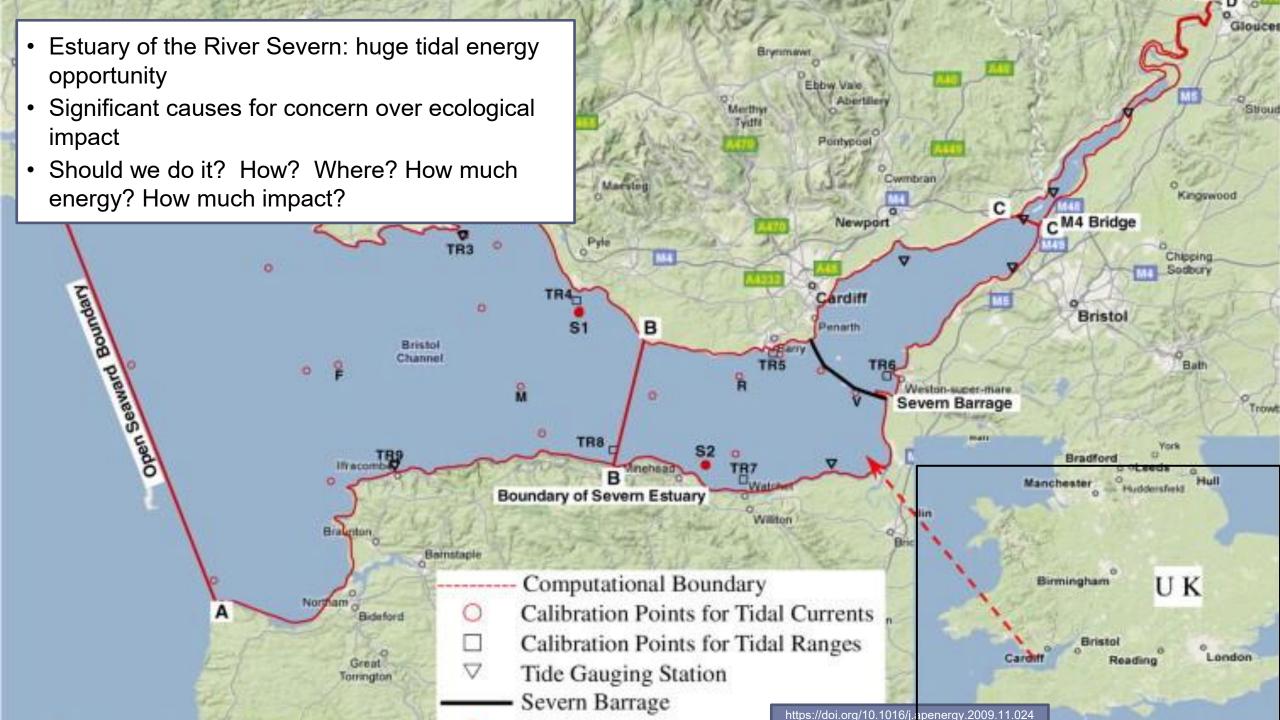
Firedrake: compiler architecture

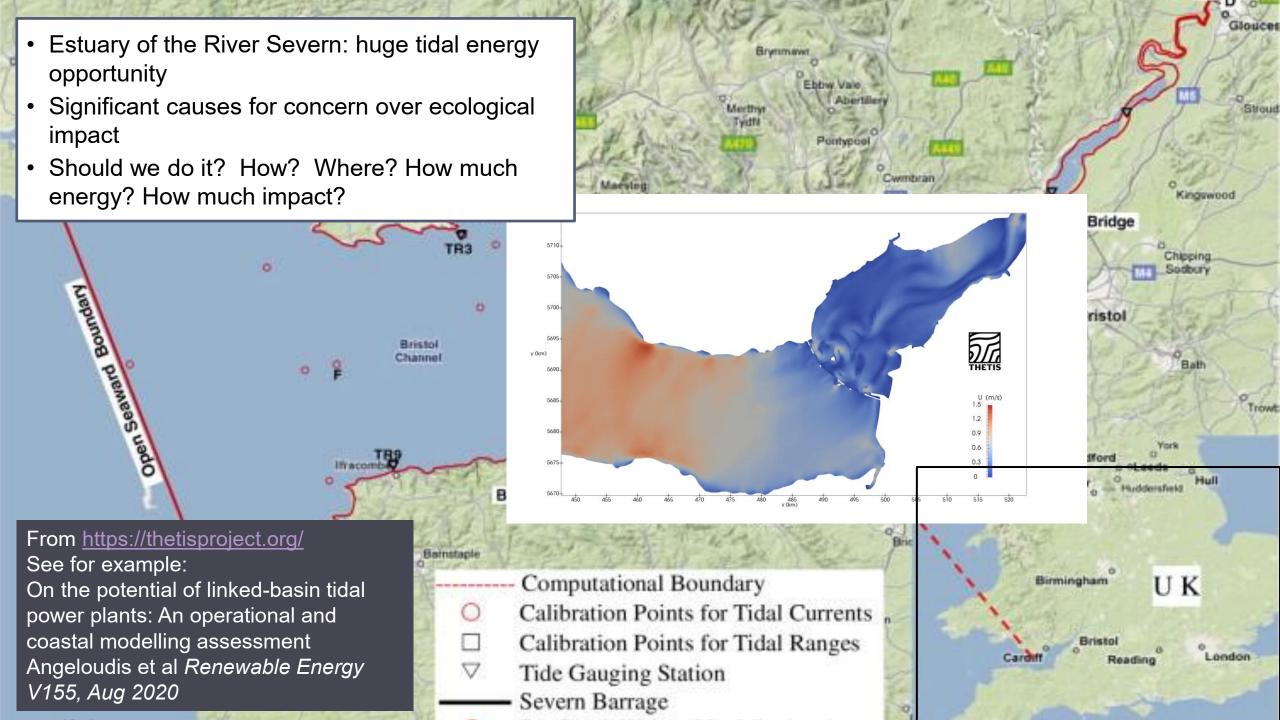


development

- Sequence of intermediate representations
- **100% Python**
- Runtime code generation, code-caching

Rathgeber, Ham, Mitchell et al, ACM TOMS 2016, Tianjiao Sun et al https://arxiv.org/pdf/1903.08243.pdf







Power tools for performance programming

By capturing domain-specific representation....

- We can deliver domain-specific optimisations
- We collect and automate all the performance techniques that are known for a family of problems
- If we get it right.... we get
 - Productivity by generating low-level code from a high-level specification
 - Performance by automating optimisations
 - Performance portability with multiple back-ends







Devito: Symbolic Finite Difference Computation

Devito is a domain-specific Language (DSL) and code generation framework for the design of highly optimised finite difference kernels for use in inversion methods. Devito utilises SymPy to allow the definition of operators from high-level symbolic equations and generates optimised and automatically tuned code specific to a given target

Symbolic computation is a powerful tool that allows users to:

architecture.

- Build complex solvers from only a few lines of high-level code
- Use automated performance optimisation for generated code
- · Adjust stencil discretisation at runtime as required
- (Re-)development of solver code in hours rather than months



https://www.devitoproject.org/

Luporini, F., Louboutin, M., Lange, M., Kukreja, N., Witte, P., Hückelheim, J., Yount, C., Kelly, P.H., Herrmann, F.J. & Gorman, G.J. Architecture and performance of Devito, a system for automated stencil computation. ACM Transactions on Mathematical Software (TOMS), 46(1) 2020



Devito: applications

- Devito automates the finite difference method for solving PDEs
- Widely used for fluid dynamics, wave propagation
- Devito is mostly used to solve inversion problems
 - Use automatic differentiation of the solver
 - To solve for the conditions that explain the observations
 - "Full Waveform Inversion" (FWI)
- Seismic inversion
 - Understand geological structures from reflected sound waves
- Ultrasound imaging of the brain
 - Diagnose brain injuries from ultrasound transmission

Devito applications largely driven by seismic inversion

- emitting transducer
- receiving transducers

Devito DSL compiler automates pathway from PDE to high-performance code

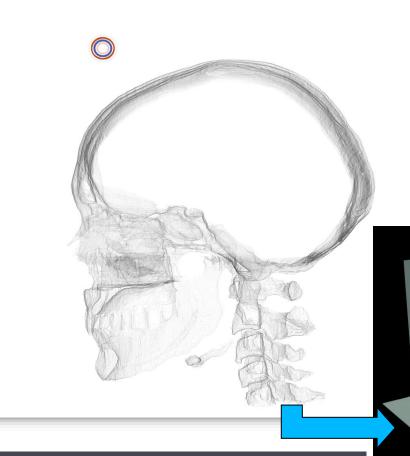
To derive forward wave propagator

And reverse adjoint wave propagator

To compute gradient

With which to correct brain model

Ultrasound imaging of the brain through the skull



Cueto, C., Guasch, L., Luporini, F., Bates, O., Strong, G., Agudo, O.C., Cudeiro, J., Kelly, P., Gorman, G. and Tang, M.X., 2022, April. Tomographic ultrasound modelling and imaging with Stride and Devito. In Medical Imaging 2022: Ultrasonic Imaging and Tomography (p. PC1203805).

vascularity derived from FWI image

Devito: example

```
# Define the wavefield from model setup.
u = TimeFunc(time_order=2, space_order=2)
# Write down the acoustic wave PDE:
pde = model.m*u.dt2 - u.laplace + model.damp*u.dt
# Solve by time-marching:
stencil = Eq(u.forward, solve(pde, u.forward))
# Define source injection and receiver:
src_term = src.inject(field=u.forward, pr=src*dt**2/model.m)
rec_term = rec.interpolate(expr=u.forward)
# Generate code for the timestepping operator:
op = Operator([stencil] + src_term + rec_term,
```

Acoustic wave equation, with damping:

$$egin{align} mrac{d^2u(x,t)}{dt^2}-
abla^2u(x,t)+\etarac{du(x,t)}{dt}=q ext{ in }\Omega\ u(.\,,0)=0\ rac{du(x,t)}{dt}|_{t=0}=0 \end{aligned}$$

We inject initial sound wave at source point, and monitor the signal at a receiver.

We derive and generate the stencil operator code, then run it a specified number of timesteps

Run code (MPI+GPU), to yield receiver values: op(time=time_range.num-1, dt=model.critical_dt)

subs=model.spacing_map)

Slightly simplified from:

https://slimgroup.github.io/Devito-Examples/tutorials/01_modelling/

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# Generate code for the timestepping operator:
op = Operator([stencil] + src_term + rec_term,
              subs=model.spacing_map)
# Run code (MPI+GPU), to yield receiver values:
op(time=time_range.num-1, dt=model.critical_dt)
```

Acoustic wave equation, with damping:

$$mrac{d^2u(x,t)}{dt^2}-
abla^2u(x,t)+\etarac{du(x,t)}{dt}=q ext{ in }\Omega \ u(.\,,0)=0 \ rac{du(x,t)}{dt}|_{t=0}=0$$

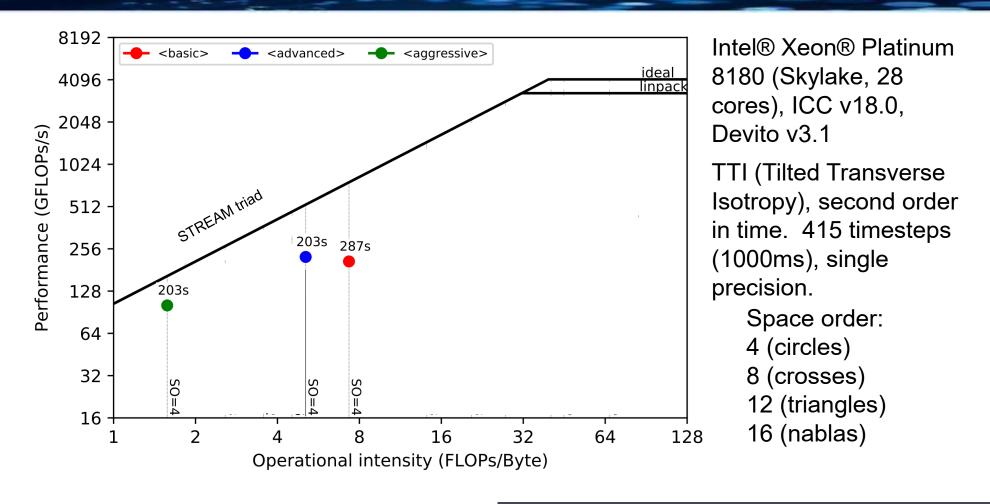
```
# No-MPI
$ python myscript.py
# With-MPI (2 ranks)
$ DEVITO_MPI=basic mpirun -n 2 python myscript.py
# MPI + GPU ready
# ...add DEVITO_PLATFORM=nvidia DEVITO_COMPILER=nvc
```

Code at this basic level of abstraction is in production, at scale, running at multiple petaflops 24/7

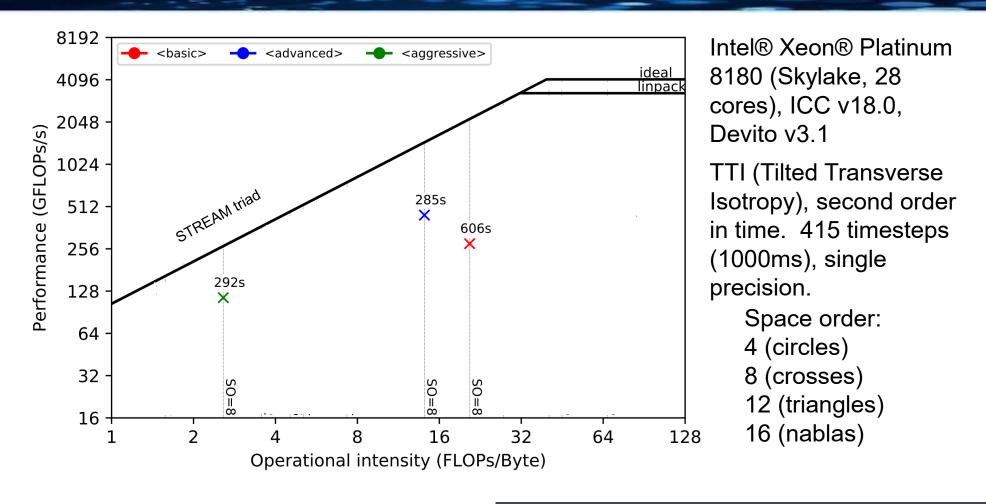
Slightly simplified from:

https://slimgroup.github.io/Devito-Examples/tutorials/01_modelling/

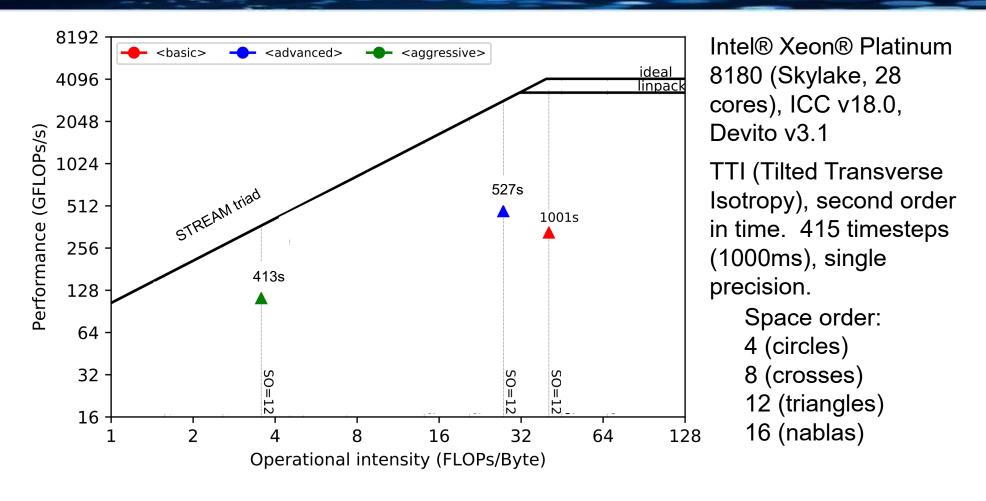




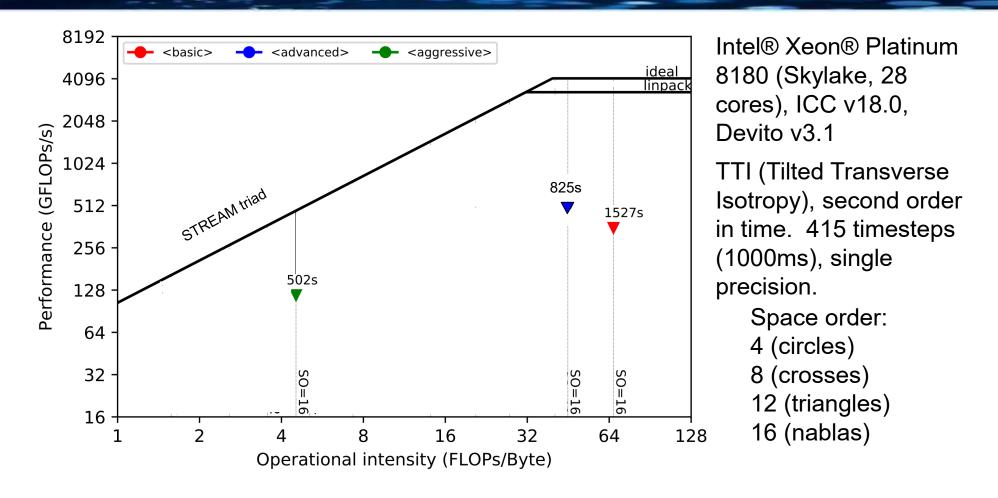




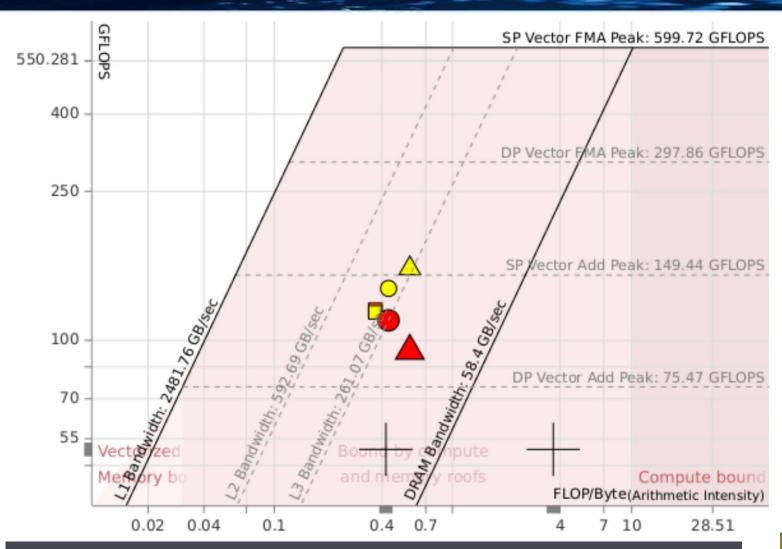








Devito: tiling-in-time



George Bisbas, et al. Temporal blocking of finite-difference stencil operators with sparse "off-the-grid" sources. IPDPS 21 iarXiv:2010.10248

Single-socket 8-core Intel Broadwell E5-2673 v4 CPUs with AVX2, L1 (32KB), L2 (256KB) private to each core, 50MB shared L3 (Ubuntu 18.04.4, Devito v4.2.3)

Isotropic acoustic model, secondorder in time, single-precision

Space order:

4 (triangles),



8 (circles), and



12 (squares).



Red markers show the performance of spatially blocked vectorized kernels

Yellow markers show spatial and temporal blocking using autotuned tile parameters.

Fixing the DSL ecosystem

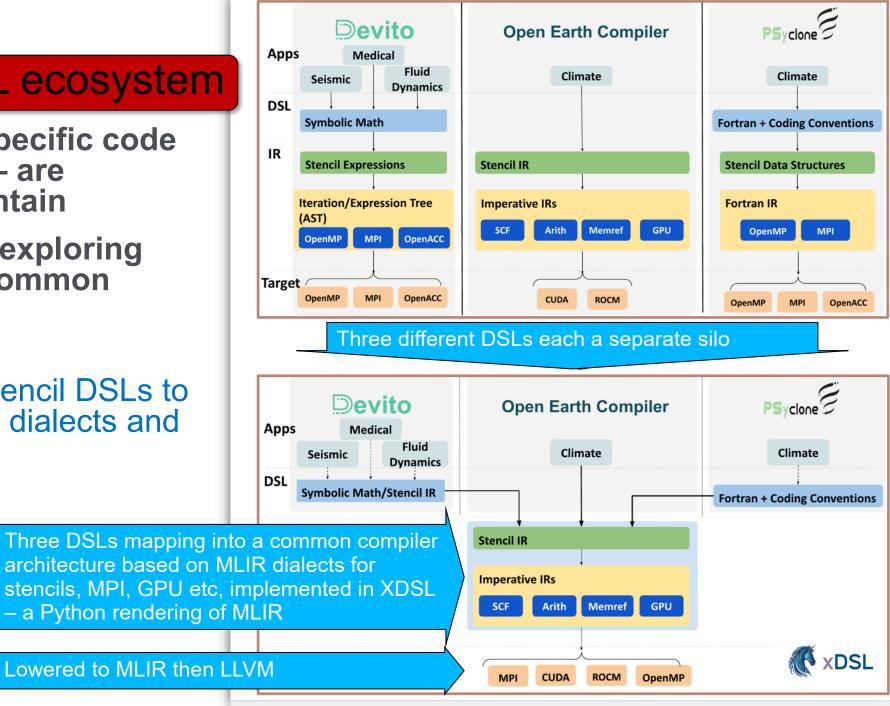
DSLs – domain-specific code generation tools - are expensive to maintain

So we have been exploring how to build on common infrastructure

In this work we restructure three stencil DSLs to use common MLIR dialects and transformations

https://xdsl.dev/

A shared compilation stack for distributed-memory parallelism in stencil DSLs Bisbas et al ASPLOS 2024









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Report repository

Releases 26

- O v0.17 (Latest on Feb 1
- + 25 releases

Contributors 54

























xDSL: A Python-native SSA Compiler Framework

codecov

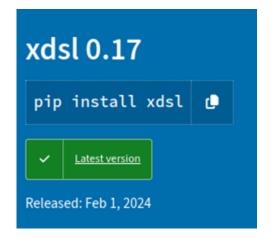
chat on zulip

pypi package 0.17

√ SSA-based IRs

CI - Python-based Testing passing

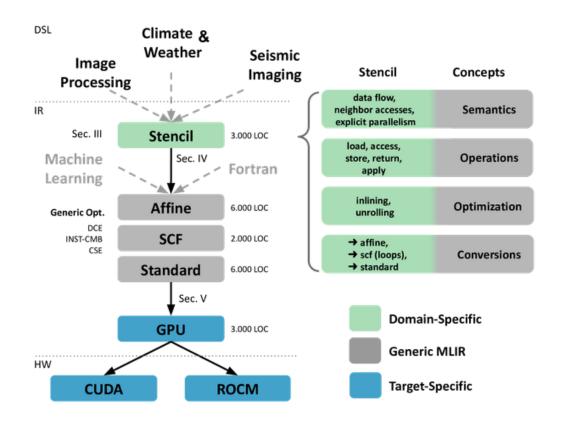
- √ SSA + regions concept
- √ Mix predefined IRs
- ✓ Add custom IRs
- √ Connect with MLIR/LLVM
- **✓** Benefit from Python's productivity 37
- **✓** Open-source/CI/CD/codecov
- √ Active contributor community
- **√** Join us on https://xdsl.zulipchat.com/



Fehr, Weber, Ulmann, Lopoukhine, Lücke, Degioanni, Vasiladiotis, Steuwer, and Grosser. 2025. XDSL: Sidekick Compilation for SSA-Based Compilers. CGO'25

The Open Earth Compiler: the 'stencil' dialect

38



```
%source = stencil.load(%114) : (!field<[0,128]xf64>)
                                    -> !temp<?xf64>
      %out = stencil.apply(%arg = %source : !temp<?xf64>)
3
                         -> !temp<?xf64> {
        %l = stencil.access %arg[-1] : f64
        %c = stencil.access %arg[0] : f64
        %r = stencil.access %arg[1] : f64
        // %v = %1 + %r - 2.0 * %c
        stencil return %v : f64
10
11
      stencil. store %out to %target([1]:[127])
12
            : !temp<?xf64> to !field<[0.128]xf64>
13
```

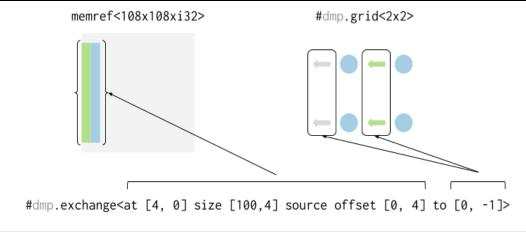
Listing 1. Example MLIR for 1-dimensional 3-point Jacobi stencil.

- ✓ Updated, ported to xDSL
- √ Extended to multi-node

Gysi et.al, Domain-Specific Multi-Level IR Rewriting for GPU: The Open Earth Compiler for GPU-accelerated Climate (2021), ACM TACO https://github.com/spcl/open-earth-compiler/

The 'dmp' and 'mpi' dialects

Listing 2. A high-level declarative expression of a data subsection exchange from some buffer.



- √ High-level halo exchanges
- ✓ Describe communication patterns
- ✓ Rectangular data subsections

'mpi' Dialect

This dialect models the Message Passing Interface (MPI), version 4.0. It is meant to serve as an interfacing dialect that is targeted by higher-level dialects. The MPI dialect itself can be lowered to multiple MPI implementations and hide differences in ABI. The dialect models the functions of the MPI specification as close to 1:1 as possible while preserving SSA value semantics where it makes sense, and uses memref types instead of bare pointers.

This dialect is under active development, and while stability is an eventual goal, it is not guaranteed at this juncture. Given the early state, it is recommended to inquire further prior to using this dialect.

For an in-depth documentation of the MPI library interface, please refer to official documentation such as the OpenMPI online documentation.

- Operations
 - mpi.comm rank (mpi::CommRankOp)
 - · mpi.error class (mpi::ErrorClassOp)
 - o mpi.finalize (mpi::FinalizeOp)
 - mpi.init (mpi::InitOp)
 - mpi.recv_(mpi::RecvOp)
 - mpi.retval check (mpi::RetvalCheckOp)
 - mpi.send (mpi::SendOp)
- Attributes
 - MPI_ErrorClassEnumAttr
- Types
 - RetvalType
 - √ Message-passing IR
 - √ Lowered to MPI library calls
 - ✓ Upstreamed to MLIR!

Stencil level IR Global to Local DMP level IR DMP to MPI MPI level IR %rank = mpi.comm_rank : i32 %ref = builtin.unrealized_conversion_cast %114 : // First swap communication calls !field<[0,64]xf64> to memref<64xf62> %dest = arith.add %rank, %minus_one : i32 %source = stencil.load(%114) : (!field<[0,128]xf64>) dmp.swap(%ref) { %is_in_bounds = arith.cmpi sge, %dest, %zero -> !temp<?xf64> "grid" = #dmp.grid<2>. scf.if %is_in_bounds { %out = stencil.apply(%arg = %source : !temp<?xf64>) "swaps" = [%view = memref.subview %ref[0][1][1] : memref<64xf64> -> !temp<?xf64> { #dmp.exchange<at [0] size [1] to memref<1xf64> %1 = stencil.access %arg[-1] : f64 source offset [1] to [-1]>, // copy data into send buffer and set up communication %c = stencil.access %arg[0] : f64 #dmp.exchange<at [64] size [1]</pre> // (omitted for clarity) %r = stencil.access %arg[1] : f64 source offset [-1] to [1]> mpi.isend %sptr, %count, %dtype, %dest, %tag, %send_req // %v = %1 + %r - 2.0 * %cmpi.irecv %rptr, %count, %dtype, %dest, %tag, %recv_req stencil.return %v : f64 }: (memref<64xf64>) -> () Exchanged %source = stencil.load(%114) ... // Second swap %out = stencil.apply(%source) ... stencil.store %out to %target([1]:[127]) // ... stencil.store %out to %target([1]:[64]) mpi.waitall %requests, %four // synchronization barrier // First swap copy back Local domain scf.if %is_in_bounds { %view = memref.subview %ref[1][1][1] : memref<64xf64> to memref<1xf64> memref.copy %recv_buffer_1 to %view 0 1 2 0 1 2 62 63 64 62 63 64 127 // Second swap copy back Global Domain Local Domains with halo exchanges highlighted // Lowered stencil comes here

Colours highlight data being operated on, shape and halo information, and communication-related information. This shows how we can enrich the IR with relevant information to perform efficient rewrites at every level of abstraction.

Stencil level IR Global to Local %source = stencil.load(%114) : (!field<[0,128]xf64>) -> !temp<?xf64> %out = stencil.apply(%arg = %source : !temp<?xf64>) -> !temp<?xf64> { %1 = stencil.access %arg[-1] : f64 %c = stencil.access %arg[0] : f64 %r = stencil.access %arg[1] : f64 // %v = %1 + %r - 2.0 * %cstencil.return %v : f64 stencil.store %out to %target([1]:[127]) 127 Global Domain

```
DMP level IR
```

DMP to MPI

MPI level IR

```
n.unrealized_conversion_cast %114
 0.641xf64> to memref<64xf62>
 mp.grid<2>,
hange<at [0] size [1]
      source offset [1] to [-1]>,
nange<at [64] size [1]
      source offset [-1] to [1]>
34xf64>) -> ()
                           Exchanged
 ncil.load(%114) ...
                                area
  .apply(%source) ....
  %out to %target([1]:[64])
            Local domain
            0 1 2
 2 63 64
                             62 63 64
ains with halo exchanges highlighted
```

```
%rank = mpi.comm_rank : i32
// First swap communication calls
%dest = arith.add %rank, %minus_one : i32
%is_in_bounds = arith.cmpi sge, %dest, %zero
scf.if %is_in_bounds {
  %view = memref.subview %ref[0][1][1] : memref<64xf64>
                                       to memref<1xf64>
  // copy data into send buffer and set up communication
  // (omitted for clarity)
  mpi.isend %sptr, %count, %dtype, %dest, %tag, %send_req
  mpi.irecv %rptr, %count, %dtype, %dest, %tag, %recv_req
// Second swap
// ...
mpi.waitall %requests, %four // synchronization barrier
// First swap copy back
scf.if %is_in_bounds {
  %view = memref.subview %ref[1][1][1] : memref<64xf64>
                                       to memref<1xf64>
  memref.copy %recv_buffer_1 to %view
// Second swap copy back
// Lowered stencil comes here
```

Colours highlight data being operated on, shape and halo information, and communication-related information. This shows how we can enrich the IR with relevant information to perform efficient rewrites at every level of abstraction.

Stencil level IR Global to Local DMP level IR %ref = builtin.unrealized_conversion_cast %114 : %source = stencil.load(%114) : (!fie !field<[0,64]xf64> to memref<64xf62> %out = stencil.apply(%arg = %source dmp.swap(%ref) { -> !temp<?xf64> { "grid" = #dmp.grid<2>, %1 = stencil.access %arg[-1] : f64 "swaps" = [%c = stencil.access %arg[0] : f64 #dmp.exchange<at [0] size [1] %r = stencil.access %arg[1] : f64 // %v = %1 + %r - 2.0 * %csource offset [1] to [-1]>, stencil.return %v : f64 #dmp.exchange<at [64] size [1]</pre> source offset [-1] to [1]> stencil.store %out to %target([1]:[1] } : (memref<64xf64>) -> () %source = stencil.load(%114) ... %out = stencil.apply(%source) ... stencil.store %out to %target([1]:[64]) Local domain Global Domain 0 1 2 62 63 64

```
DMP to MPI
                              MPI level IR
```

Exchanged

area

62 63 64

Local Domains with halo exchanges highlighted

```
%rank = mpi.comm_rank : i32
// First swap communication calls
%dest = arith.add %rank, %minus_one : i32
%is_in_bounds = arith.cmpi sge, %dest, %zero
scf.if %is_in_bounds {
  %view = memref.subview %ref[0][1][1] : memref<64xf64>
                                       to memref<1xf64>
  // copy data into send buffer and set up communication
  // (omitted for clarity)
  mpi.isend %sptr, %count, %dtype, %dest, %tag, %send_req
  mpi.irecv %rptr, %count, %dtype, %dest, %tag, %recv_req
// Second swap
// ...
mpi.waitall %requests, %four // synchronization barrier
// First swap copy back
scf.if %is_in_bounds {
  %view = memref.subview %ref[1][1][1] : memref<64xf64>
                                       to memref<1xf64>
  memref.copy %recv_buffer_1 to %view
// Second swap copy back
// Lowered stencil comes here
```

Colours highlight data being operated on, shape and halo information, and communication-related information. This shows how we can enrich the IR with relevant information to perform efficient rewrites at every level of abstraction.

Global to Local

%source = stencil.load(%114) : (!field<[0,128]xf64>) -> !temp<?xf64> %out = stencil.apply(%arg = %source : !temp<?xf64>) -> !temp<?xf64> { %1 = stencil.access %arg[-1] : f64 %c = stencil.access %arg[0] : f64 %r = stencil.access %arg[1] : f64 // %v = %1 + %r - 2.0 * %cstencil.return %v : f64 stencil.store %out to %target([1]:[127]) 127 Global Domain

Stencil level IR

```
%ref = builtin.unrealized
        !field<[0.64]xf64> to
   dmp.swap(%ref) {
     "grid" = #dmp.grid<2>,
     "swaps" = [
       #dmp.exchange<at [0]
                      source
       #dmp.exchange<at [64]
                      source
   }: (memref<64xf64>) -> (
   %source = stencil.load(%1
   %out = stencil.apply(%sour
   stencil.store %out to %tar
0 1 2
               62 63 64
       Local Domains with halo exc
```

DMP level IR

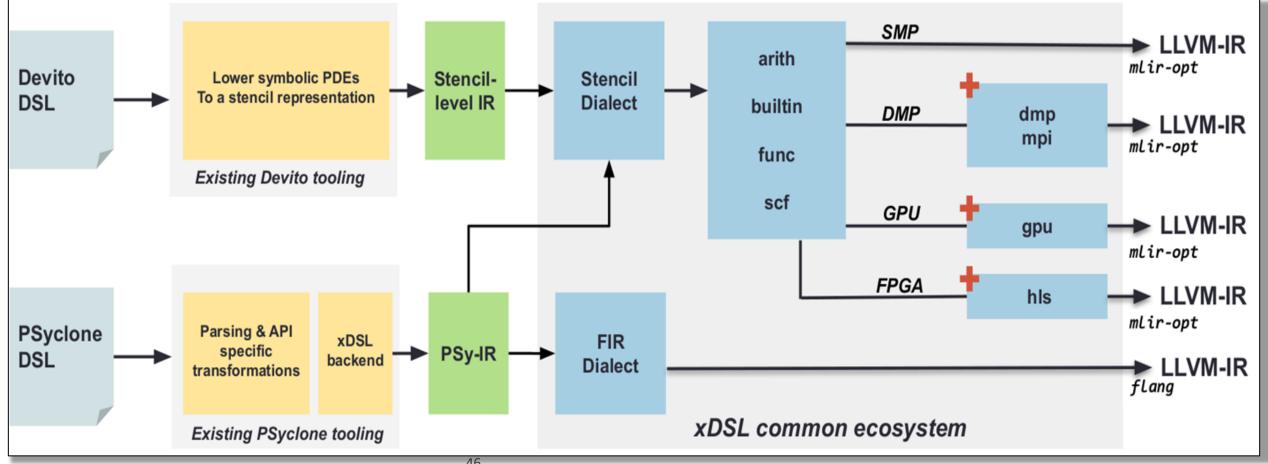
DMP to MPI

Colours highlight data being operated on, shape and halo information, and communication-related information. This shows how we can enrich the IR with relevant information to perform efficient rewrites at every level of abstraction.

```
%rank = mpi.comm_rank : i32
// First swap communication calls
%dest = arith.add %rank, %minus_one : i32
%is_in_bounds = arith.cmpi sge, %dest, %zero
scf.if %is in bounds {
  %view = memref.subview %ref[0][1][1] : memref<64xf64>
                                       to memref<1xf64>
  // copy data into send buffer and set up communication
  // (omitted for clarity)
  mpi.isend %sptr, %count, %dtype, %dest, %tag, %send_req
  mpi.irecv %rptr, %count, %dtype, %dest, %tag, %recv_req
// Second swap
mpi.waitall %requests, %four // synchronization barrier
// First swap copy back
scf.if %is_in_bounds {
  %view = memref.subview %ref[1][1][1] : memref<64xf64>
                                       to memref<1xf64>
  memref.copy %recv_buffer_1 to %view
  Second swap copy back
// Lowered stencil comes here
```

MPI level IR

The shared compilation stack for DMP in stencil DSLs

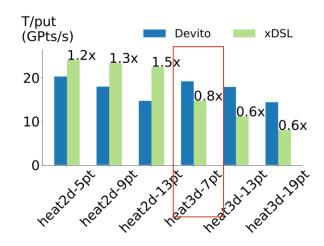


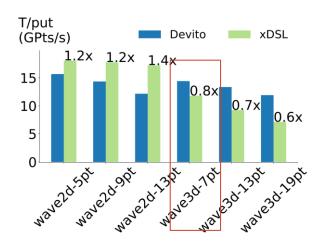
- ✓ Unlocked optimizations
- ✓ Unlocked multi-node CPU
- ✓ Unlocked other backends (FPGA, CUDA)
- ✓ Competitive or better performance with an order of 1000s LoC saved!

Gysi et.al, Domain-Specific Multi-Level IR Rewriting for GPU: The Open Earth Compiler for GPU-accelerated Climate (2021), ACM TACO https://github.com/spcl/open-earth-compiler/

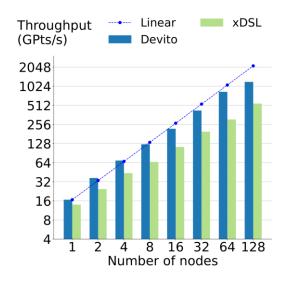
Performance evaluation: Devito

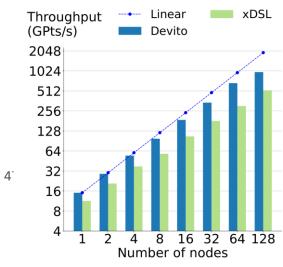




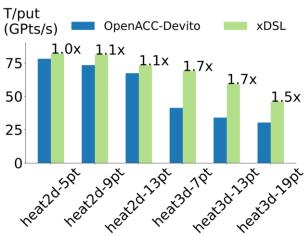


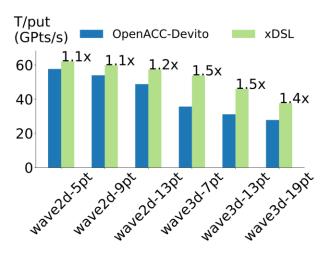
Single-node AMD EPYC 7742, 8 MPI ranks x 16 OpenMP threads, 16384^2 (2D) and 1024^3 (3D)





Heat (top) and wave (bottom), 3D-7pt, multinode strong scaling up to 128 nodes, total of 16384 cores.





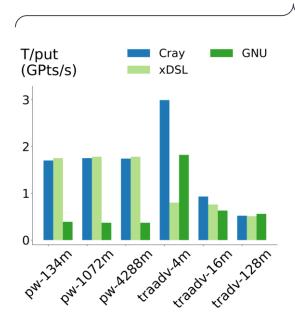
xDSL adds support for CUDA, outperforming Devito's OSS support for OpenACC, running on V100-SXM2-16GB (Volta).

Performance evaluation: PSyclone



Piacsek and Williams (PW) advection and NEMO tracer advection (traadv) kernels

T/put



(GPts/s) PSyclone xDSL

x24.14

x14.60x11.01

x0.62 x0.83 x0.95

pm. 33m

pm. 134m

pm. 134m

kraadv. Am

kraadv. 128m

kraadv. 128m

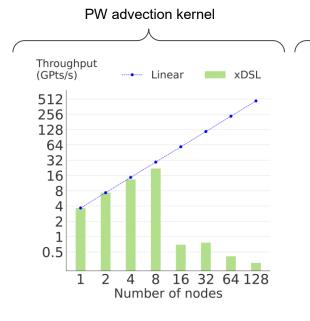
kraadv. 128m

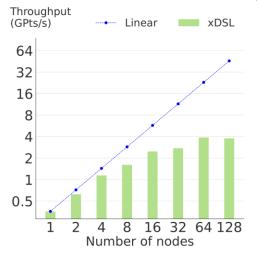
Single-node AMD EPYC 7742 throughput, PSyclone target code compiled with Cray and GNU compilers against xDSL-PSyclone.

xDSL-PSyclone code matches Cray code and significantly outperforms GNU code for PW advection.

Single-node Cirrus NVIDIA Tesla V100-SXM2-16GB throughput, PSyclone NVIDIA GPU code against xDSL-PSyclone GPU code.

xDSL-PSyclone significantly outperforms PSyclone NVIDIA for PW advection due to data allocation approach (explicit device memory allocation for xDSL-PSyclone).





traady kernel

Multi-node strong scaling of problem size [256,256,128], scaling up to 128 nodes, total of 16384 cores.

Suffers scaling effects at 8 nodes due to small global problem size.

Multi-node strong scaling of problem size [512,512,128], scaling up to 128 nodes, total of 16384 cores.

2D decomposition strategy limits strong scaling.

Conclusions

- Productivity by generating low-level code from a high-level specification
- Performance by automating optimisations
- Performance portability with multiple back-ends
- Domain-specific compiler design is all about designing representations that make hard problems easy
- The grand project is to build common compiler infrastructure that spans different domains
- DSL compilers exploit data structures
- DSL compilers exploit computation
- Redundancy
- Locality
- Parallelism

- You can actually have it all, today
- Tensor contractions
- Access/execute
- With MLIR
- Via XDSL
- Meshes
- Metadata
- Taming pointers
- Composition
- Adaptivity

- Thank you to:
 - EPSRC
 - EP/Y020499/1 On-Sensor Computer Vision
 - EP/W026066/1 SysGenX: Composable software generation for system-level simulation at Exascale
 - EP/W007789/1 Efficient Cross-Domain DSL Development for Exascale
 - EP/V001493/1 Gen X: ExCALIBUR working group on Exascale continuum mechanics through code generation
 - EP/R029423/1 PRISM: Platform for Research In Simulation Methods
 - EP/P010040/1 Application Customisation: Enhancing Design Quality and Developer Productivity
 - EP/K008730/1 PAMELA: a Panoramic Approach to the Many-CorE LAndscape from end-user to end-device: a holistic game-changing approach
 - EP/I00677X/1 Multi-layered abstractions for PDEs